Verified Secure Routing

Peter Müller

Joint work with the Verified Scion Team at ETH





- Internet is a network of Autonomous Systems (AS)
- Each AS is a network of routers run by an institution
- Routes between AS are discovered using Border Gateway Protocol (BGP)
- Based on trust, for instance, any AS can announce any IP address range

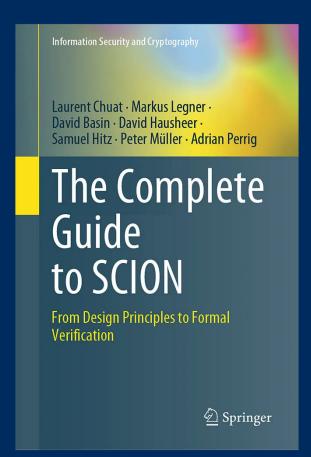
There are numerous ways to attack Internet routing



- In 2013, Ukrainian ISP announced route prefixes to British Telecom AS
- Traffic of some UK customers was redirected to Ukraine, including UK's Atomic Weapons Establishment
- Senders have no control over the taken routes
- Routers on path can read and modify data

Scion Internet Architecture

- Scion is a new architecture for inter-domain routing
 - Path control, e.g., geofencing
 - Multipath communication
 - DDoS protection
- Research and commercial deployments

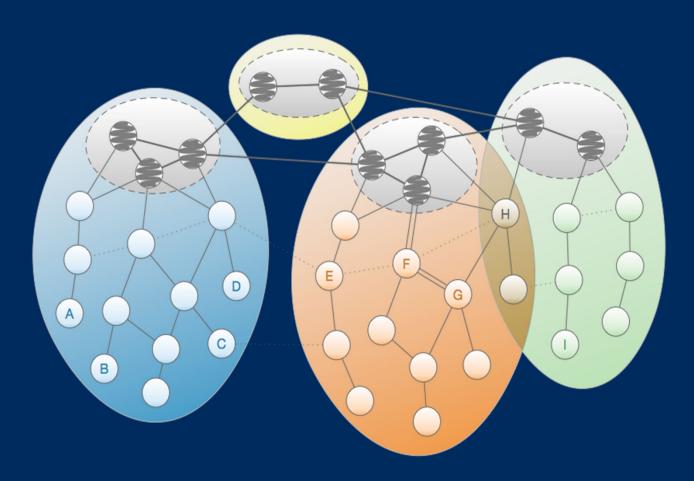




Formal end-to-end verification of security and correctness

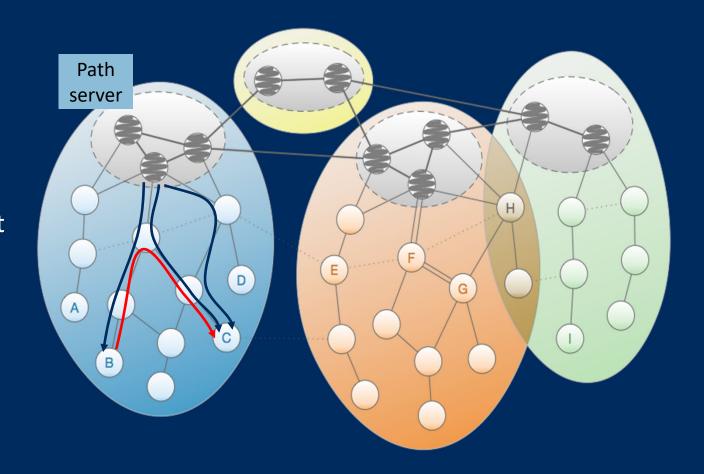
Isolation Domains

ASes are organized into isolation domains with independent control planes and root of trust



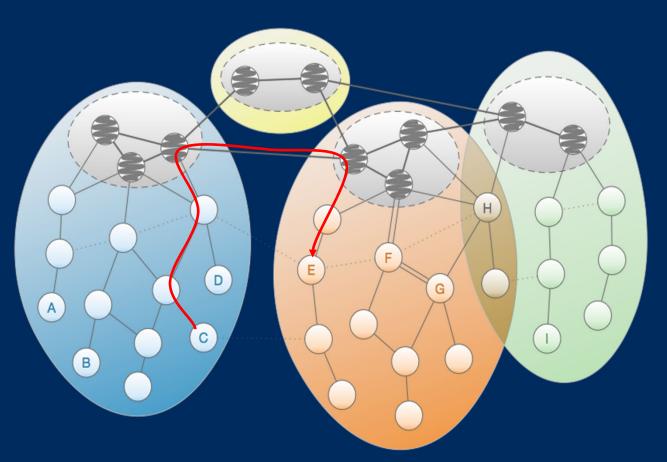
Scion Routing

- Path exploration
 - Paths are sequences of signed hop fields
 - Each hop field carries routing information for one AS (input and output ports)
- Path registration with path server
- Path selection
 - Path is stored in packet header



Scion Forwarding

- Path is stored in packet header
- Consisting of up segment, core segment, and down segment



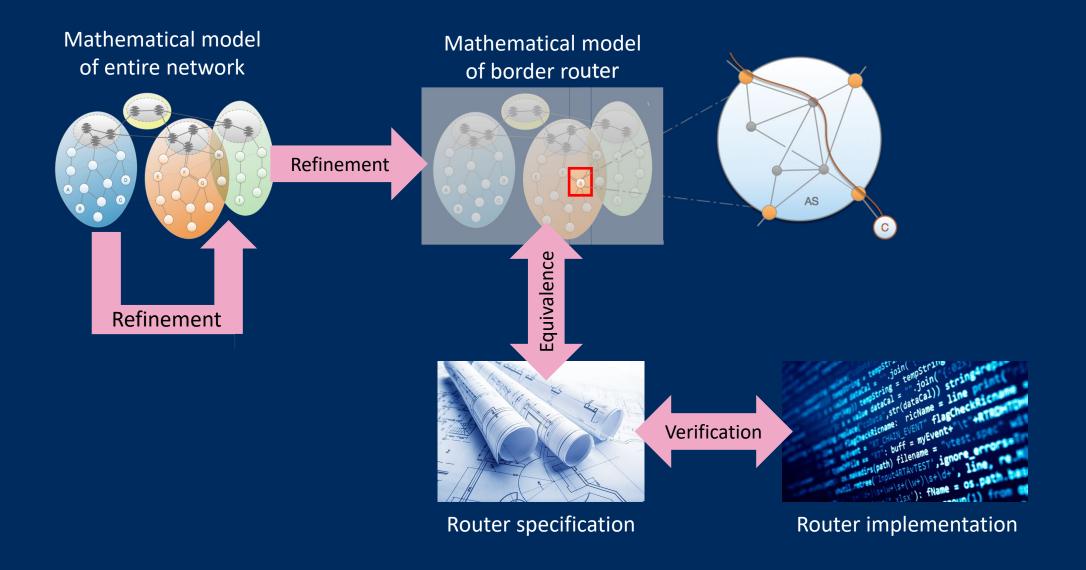
Security and Correctness

Protocol-level properties

- **Path validity**: Constructed paths are valid and reflect the routing decisions by on-path ASes.
- Path authorization: Packets travel only along previously authorized paths
- **Detectability**: An active attacker cannot hide their presence on the path

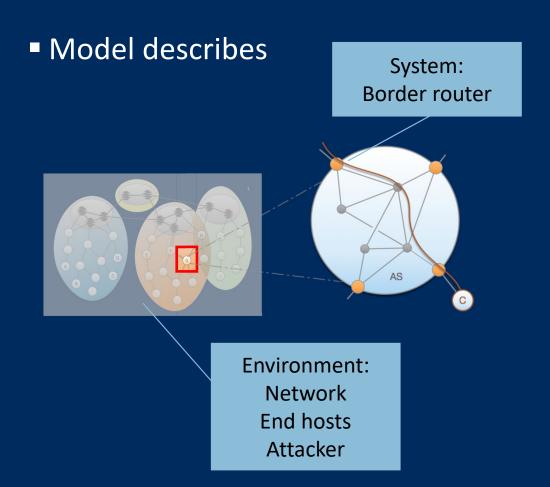
Code-level properties

- Safety: No run-time errors
- **Correctness**: Routers and servers implement protocol correctly
- **Progress**: Required I/O happens eventually
- **Secure information flow**: Code does not leak information about crypto keys



Design Model

Formalize the design model as transition system

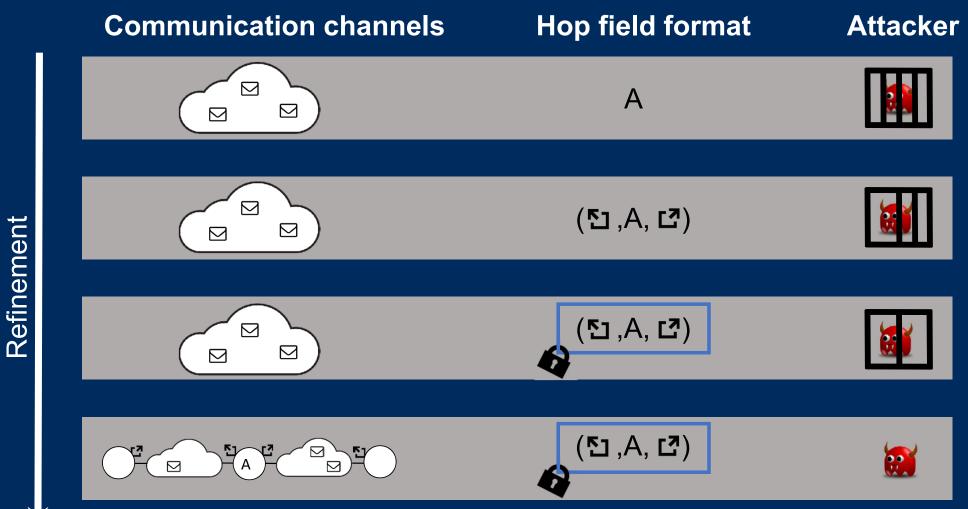


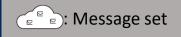
Stepwise Refinement

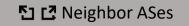
- Design model is developed by stepwise refinement
- Prove properties of most abstract model
- Each refinement
 - Incorporates additional system requirements or environment assumptions
 - Preserves properties of more-abstract system
 - Is tool-checked in Isabelle



Strategy: strengthen attacker while increasing protection of paths





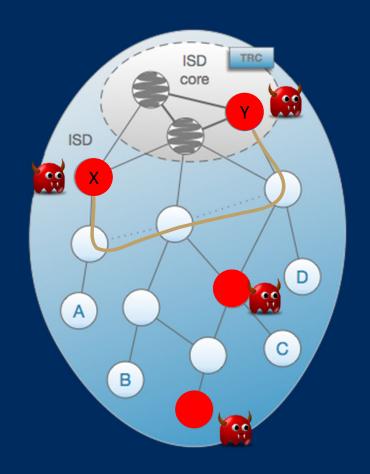




Attacker Model

Localized, colluding Dolev-Yao attacker model

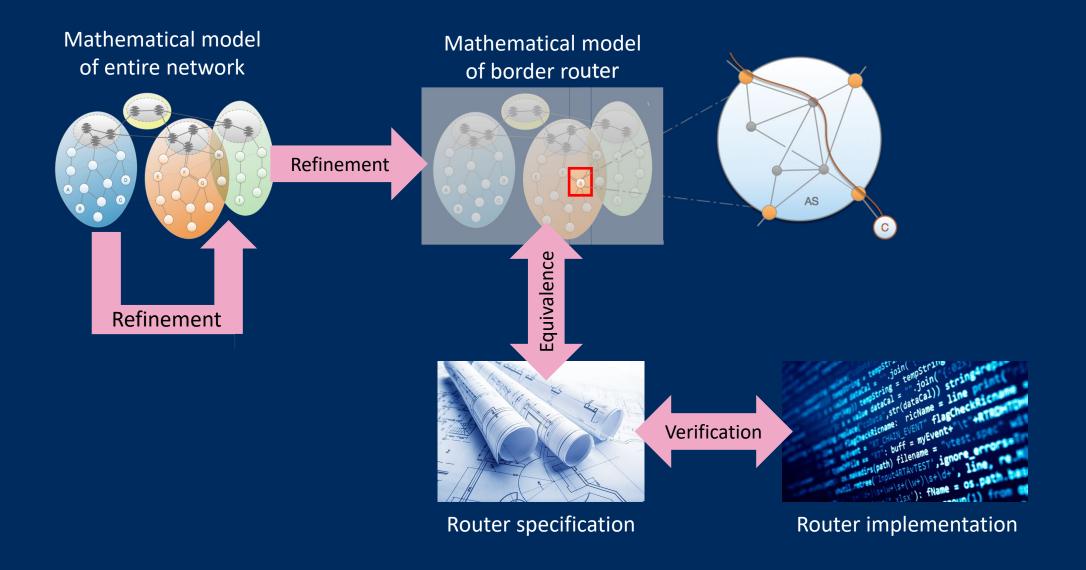
- Attacker:
 - Actively controls some ASes
 - Can observe, block, and inject messages
 - Can eavesdrop globally
- Cryptography is assumed to be perfect



Results of Protocol Verification

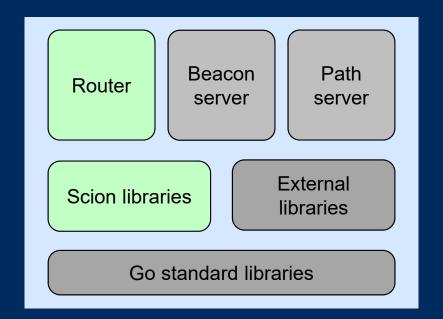
- A formal model of the network components and their environment
- Model serves as formal specification for the implementation
- Proofs of the desired properties under the assumption that each component satisfies its specification
- 16,100 lines (models and proofs)

- Improved understanding of protocols and properties
- Revealed design flaws that enabled five different security attacks
- Issues were found during modeling and formalization



Scion Implementation

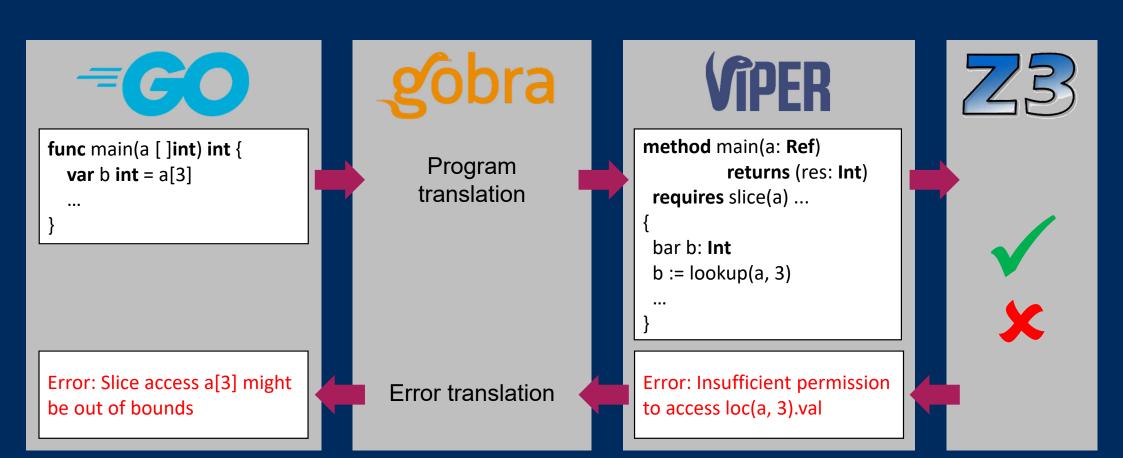
- Open-source implementation
 - 35kloc of Go (Router: 4.7kloc)
 - Uses concurrency, async, globals
- Verify safety, functional correctness, progress, secure information flow
- Assume correctness of external libraries,
 Go compiler, OS, hardware



Gobra: Sound Verification for Go

- No run-time errors
- No data races
- Functional properties
- Termination
- I/O behavior
- Secure information flow

Gobra Toolchain



Permissions

- Associate each heap location with a permission
- Permissions are held by method executions
- Access to a memory location requires permission

```
func indexOf(| [ ]int, i, val int) (res int)
  requires forall j int :: i <= j && j < len(|) ==> acc(&|[ j ])
  ensures forall j int :: i <= j && j < len(|) ==> acc(&|[ j ])
```

- Permissions can be transferred, but not duplicated or forged
- Predicates abstract over concrete permissions

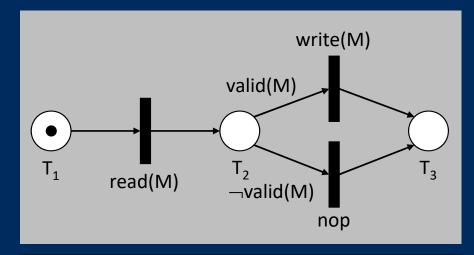
I/O Permissions

- Permissions and predicates can be used to reason about resources
- Here: permission to perform an I/O operation

func write(value string)
requires write(O(value))

I/O Behavior as Petri Nets

- Adaptation of work by Penninckx et al.
- Petri nets specify permitted I/O behavior
 - Traces of basic I/O operations
 - Sequences, parallelism, non-determinism
- Petri nets are encoded as (recursive) predicates



```
predicate router(T<sub>1</sub>) {

∀M ∃T<sub>2</sub>, T<sub>3</sub> •

readIO(T<sub>1</sub>,M,T<sub>2</sub>) *

(valid(M) ⇒ writeIO(T<sub>2</sub>,M,T<sub>3</sub>)) *

(¬ valid(M) ⇒ nop(T<sub>2</sub>,T<sub>3</sub>)) *

router(T<sub>3</sub>)
}
```

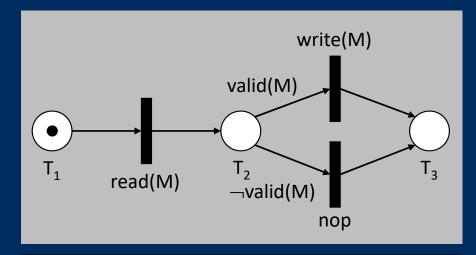
Specification of I/O Behavior

- Basic I/O operations
 - Require I/O permission
 - Require token in appropriate place
 - Advance token

```
func write(value string)
  requires token(T) * writeIO(T, value, T')
  ensures token(T')
```

Method precondition characterizes permitted I/O behavior

```
func main()
  requires token(T) * router(T)
```



```
predicate router(T<sub>1</sub>) {

∀M ∃T<sub>2</sub>, T<sub>3</sub> •

readIO(T<sub>1</sub>,M,T<sub>2</sub>) *

(valid(M) ⇒ writeIO(T<sub>2</sub>,M,T<sub>3</sub>)) *

(¬ valid(M) ⇒ nop(T<sub>2</sub>,T<sub>3</sub>)) *

router(T<sub>3</sub>)
}
```

From Design Model to I/O Specification

■ Refine design model to have an event for each basic I/O operation

I/O event write(val)

func write(value string)

Encode entire event system as recursive predicate

```
event drop(M)
guard ¬valid(M)
action
buf := buf \ { M }
```

```
predicate system(T<sub>1</sub>, state) { 
 (∀args • guard(args, state) ⇒ 
 \exists T_2 • oplO(T<sub>1</sub>,args,T<sub>2</sub>) * system(T<sub>2</sub>, state')) * 
 ...
```

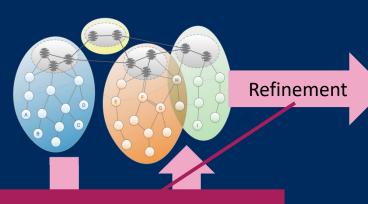
```
predicate router(T<sub>1</sub>, buf) { (∀M \bullet \neg valid(M) \Rightarrow router(T<sub>1</sub>, buf \ { M })) * ...}
```

Status of Code Verification

- Completed verification of SCION router (4,700 LoC)
 - Memory safety
 - Functional correctness
 - I/O behavior
 - Termination
- 13,400 lines of annotations (2.8 LoS per LoC)
- Required only three code changes

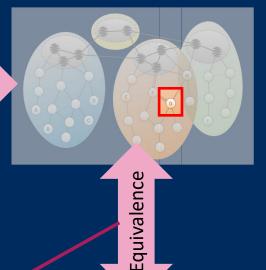
- Identified 13 confirmed issues related to memory safety, functional correctness, and I/O behavior (plus 2 performance issues)
- Despite extensive code reviews, testing, and fuzzing
- Verification is currently being introduced into the continuous integration

Mathematical model of entire network



- Developed by stepwise refinement
- Verified in Isabelle
 - Transition system encoded as I/O specification
 - Verified in Isabelle

Mathematical model of border router



- Properties encoded in permission logic
- Verified in Gobra/Viper



Verification

Router specification



Router implementation